

# ECS455: Chapter 5

## OFDM

### 5.3 Implementation: DFT and FFT



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**Thursday 13:30-14:30**

# Discrete Fourier Transform (DFT)

Transmitter produces

$$s(t) = \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} S_k \exp\left(j \frac{2\pi k}{T_s} t\right), \quad 0 \leq t < T_s$$



Sample the signal **in time domain** every  $T_s/N$  gives

sampling rate

$$= \frac{1}{T_s/N} = \frac{N}{T_s}$$

$$s[n] = s\left(n \frac{T_s}{N}\right) = \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} S_k \exp\left(j \frac{2\pi k}{T_s} n \frac{T_s}{N}\right)$$

$s(t)$  |  $t = n \frac{T_s}{N}$

$$= \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} S_k \exp\left(j \frac{2\pi kn}{N}\right) = \sqrt{N} \text{IDFT}\{S\}[n]$$

where  $\text{IDFT}\{S\}[n] = \frac{1}{N} \sum_{k=0}^{N-1} S_k \exp\left(j \frac{2\pi kn}{N}\right)$

$$S = (S_0 \quad S_1 \quad \dots \quad S_{N-1})^T$$

We can implement OFDM in the discrete domain!

# Discrete Fourier Transform (DFT)

Here, we work with  $N$ -point signals (finite-length sequence (vector) of length  $N$ ) in both time and frequency domain.

$$\mathbf{x} = \begin{pmatrix} x[0] \\ x[1] \\ \vdots \\ x[N-1] \end{pmatrix} \xrightarrow{\text{DFT}} \mathbf{X} = \begin{pmatrix} X[0] \\ X[1] \\ \vdots \\ X[N-1] \end{pmatrix}$$

$$\mathbf{X} = \text{DFT}\{\mathbf{x}\} = \Psi_N \mathbf{x}$$

To simplify the notation, we define the **DFT matrix**  $\Psi_N$ .

$$(\Psi_N)^{-1} \mathbf{X} = \text{IDFT}\{\mathbf{X}\} = \mathbf{x} \xrightleftharpoons[\text{IDFT}]{\text{DFT}} \mathbf{X} = \text{DFT}\{\mathbf{x}\} = \Psi_N \mathbf{x}$$

# DFT matrix $\Psi_N$

$$\Psi_N = \begin{bmatrix} 1 & 1 & 1 & \dots & 1 \\ 1 & \psi_N^{-1} & \psi_N^{-2} & \dots & \psi_N^{-(N-1)} \\ 1 & \psi_N^{-2} & \psi_N^{-4} & \dots & \psi_N^{-2(N-1)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \psi_N^{-(N-1)} & \psi_N^{-2(N-1)} & \dots & \psi_N^{-(N-1)(N-1)} \end{bmatrix}$$

Element on the  $p$ th row and  $q$ th column is given by

$$\psi_N^{-(p-1)(q-1)} \text{ where } \psi_N = e^{j\frac{2\pi}{N}}$$

Note that the “-1” are there because we start from row 1 and column 1 (not from row 0 and column 0).

Key Property:  $\Psi_N^{-1} = \frac{1}{N} \Psi_N^*$   $\xrightarrow{\text{complex conjugate}}$   $\frac{1}{\sqrt{N}} \Psi_N$  is a unitary matrix

$$\psi_2 = e^{j\frac{2\pi}{2}} = e^{j\pi} = -1$$

$$\Psi_2 = \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

$$\frac{1}{N} \Psi_N^* \mathbf{X} = (\Psi_N)^{-1} \mathbf{X} = \text{IDFT} \{ \mathbf{X} \} = \mathbf{x} \xrightleftharpoons[\text{IDFT}]{\text{DFT}} \mathbf{X} = \text{DFT} \{ \mathbf{x} \} = \Psi_N \mathbf{x}$$

$$\vec{x} = \begin{bmatrix} 2 \\ 5 \end{bmatrix} \quad \vec{X} = \text{DFT}(\vec{x}) = \Psi_2 \vec{x} = \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} 2 \\ 5 \end{bmatrix} = \begin{bmatrix} 7 \\ -3 \end{bmatrix}$$

$$\text{IDFT}(\vec{x}) = \frac{1}{2} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} 7 \\ -3 \end{bmatrix} = \begin{bmatrix} 2 \\ 5 \end{bmatrix}$$

# Discrete Fourier Transform (DFT)

Matrix form:

$$\frac{1}{N} \Psi_N^* \mathbf{X} = \text{IDFT} \{ \mathbf{X} \} = \mathbf{x} \begin{array}{c} \xrightarrow{\text{DFT}} \\ \xleftarrow{\text{IDFT}} \end{array} \mathbf{X} = \text{DFT} \{ \mathbf{x} \} = \Psi_N \mathbf{x}$$

Pointwise form:

$$\frac{1}{N} \sum_{k=0}^{N-1} X[k] \psi_N^{nk} = x[n] \begin{array}{c} \xrightarrow{\text{DFT}} \\ \xleftarrow{\text{IDFT}} \end{array} X[k] = \sum_{n=0}^{N-1} x[n] \psi_N^{-nk}$$

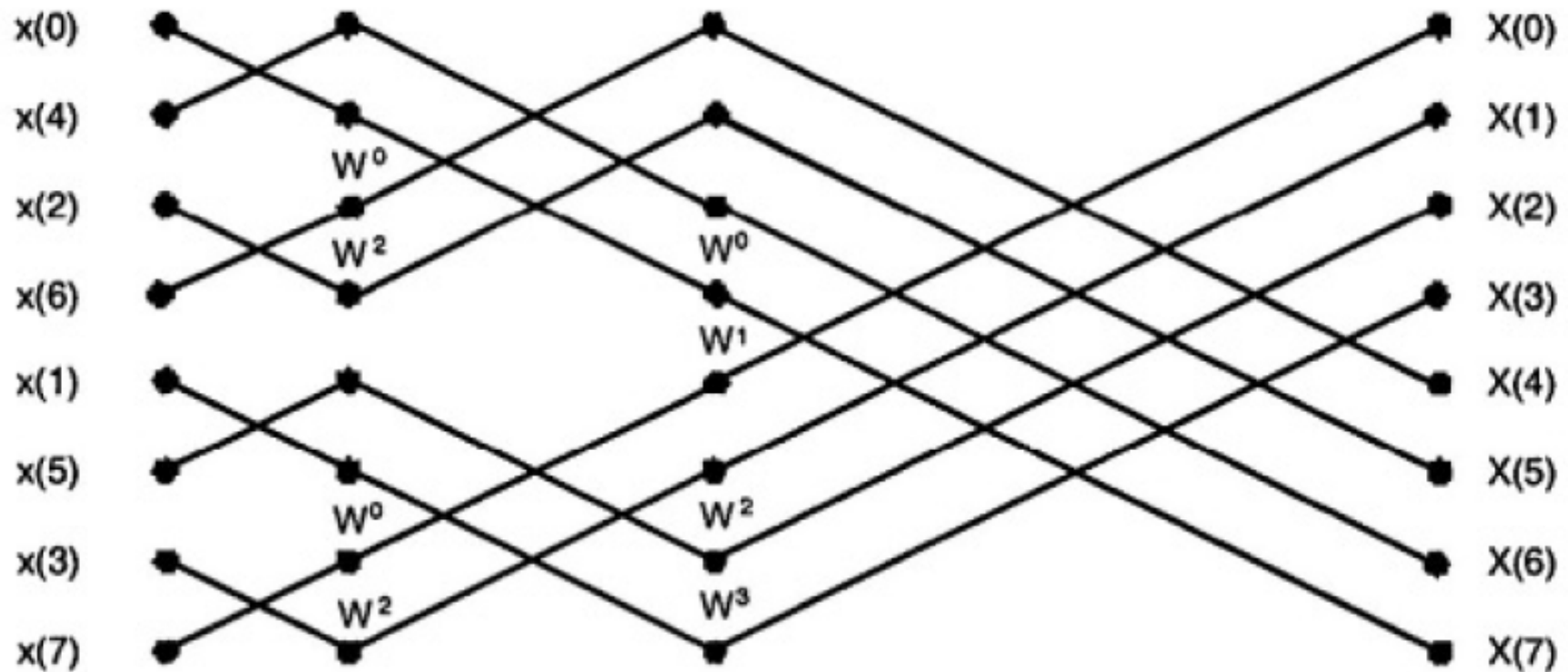
or, equivalently,

$$\frac{1}{N} \sum_{n=0}^{N-1} X[k] e^{jnk \frac{2\pi}{N}} = x[n] \begin{array}{c} \xrightarrow{\text{DFT}} \\ \xleftarrow{\text{IDFT}} \end{array} X[k] = \sum_{n=0}^{N-1} x[n] e^{-jnk \frac{2\pi}{N}}$$

Comparison with Fourier transform

$$x(t) = \int_{-\infty}^{\infty} X(f) e^{j2\pi ft} df \begin{array}{c} \xrightarrow{\mathcal{F}} \\ \xleftarrow{\mathcal{F}^{-1}} \end{array} x(f) = \int_{-\infty}^{\infty} X(t) e^{-j2\pi ft} dt$$

# Efficient Implementation: (I)FFT

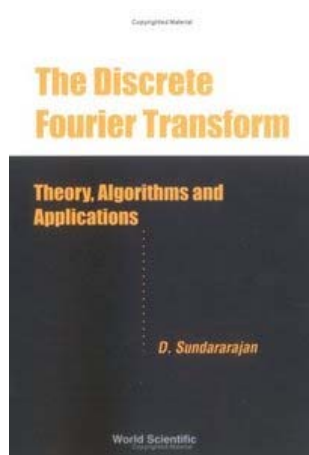


[Bahai, 2002, Fig. 2.9]

An  $N$ -point FFT requires only on the order of  $N \log N$  multiplications, rather than  $N^2$  as in a straightforward computation.

# FFT

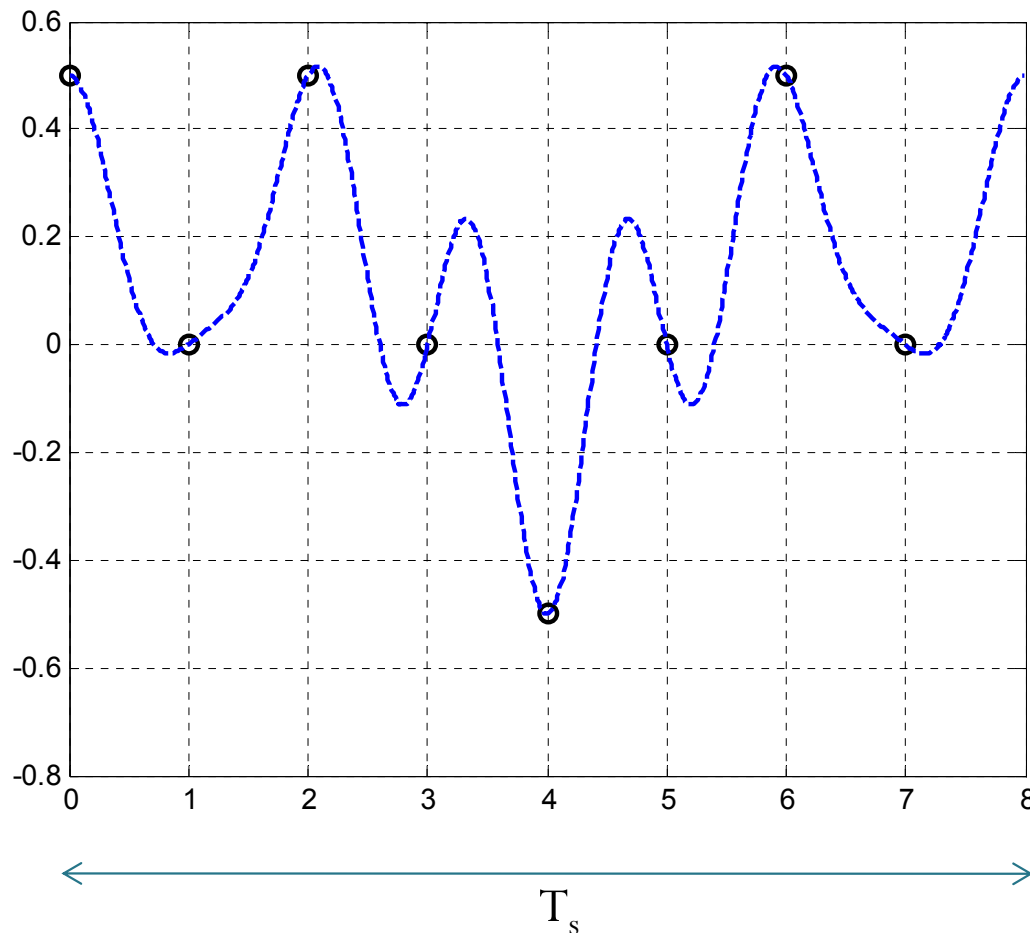
- The history of the FFT is complicated.
- As with many discoveries and inventions, it arrived before the (computer) world was ready for it.
- Usually done with  $N$  a power of two.
  - Very efficient in terms of computing time
  - Ideally suited to the binary arithmetic of digital computers.
  - Ex: From the implementation point of view it is better to have, for example, a FFT size of 1024 even if only 600 outputs are used than try to have another length for FFT between 600 and 1024.



References: E. Oran Brigham, *The Fast Fourier Transform*, Prentice-Hall, 1974.

# DFT Samples

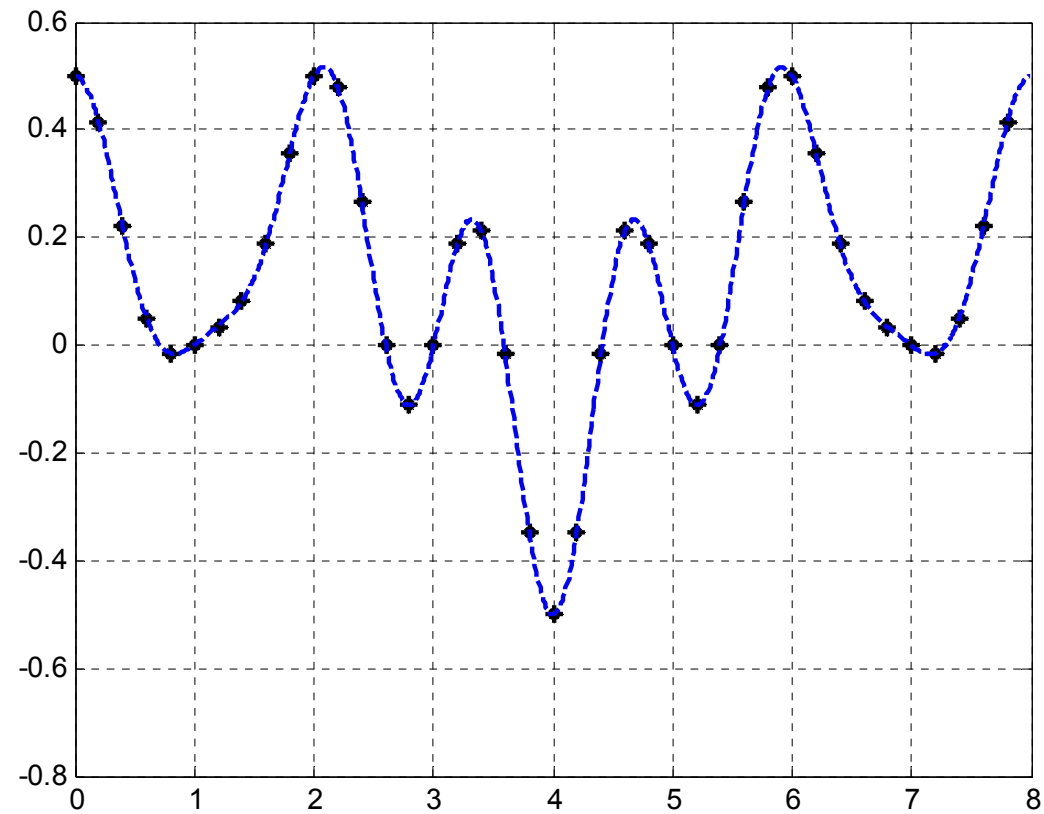
$$s(t) = \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} S_k \exp\left(j \frac{2\pi kt}{T_s}\right), \quad 0 \leq t \leq T_s$$



$$\begin{aligned} s[n] &= s\left(n \frac{T_s}{N}\right) \\ &= \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} S_k \exp\left(j \frac{2\pi kn}{N}\right) \\ &= \sqrt{N} \text{IDFT}\{S\}[n] \\ &0 \leq n < N \end{aligned}$$



# Oversampling



# Oversampling (2)

- Increase the number of sample points from  $N$  to  $LN$  on the interval  $[0, T_s]$ .
- $L$  is called the **over-sampling factor**.

$$s[n] = s\left(n \frac{T_s}{N}\right) \quad 0 \leq n < N \quad \longrightarrow \quad s^{(L)}[n] = s\left(n \frac{T_s}{LN}\right) \quad 0 \leq n < LN$$

$$\begin{aligned} s^{(L)}[n] &= \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} S_k \exp\left(j \frac{2\pi k}{T_s} n \frac{T_s}{LN}\right) = \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} S_k \exp\left(j \frac{2\pi kn}{LN}\right) \\ &= \frac{1}{\sqrt{N}} LN \left( \frac{1}{LN} \sum_{k=0}^{N-1} S_k \exp\left(j \frac{2\pi kn}{LN}\right) \right) \\ &= L\sqrt{N} \left( \frac{1}{LN} \left( \sum_{k=0}^{N-1} S_k \exp\left(j \frac{2\pi kn}{LN}\right) + \sum_{k=N}^{LN-1} 0 \exp\left(j \frac{2\pi kn}{LN}\right) \right) \right) \\ &= L\sqrt{N} \left( \frac{1}{LN} \sum_{k=0}^{LN-1} \tilde{S}_k \exp\left(j \frac{2\pi kn}{LN}\right) \right) = L\sqrt{N} \text{IDFT}\{\tilde{S}\}[n] \end{aligned}$$

Zero padding:

$$\tilde{S}_k = \begin{cases} S_k, & 0 \leq k < N \\ 0, & N \leq k < LN \end{cases}$$

Scaling

# Oversampling: Summary

$N$  points

$$s[n] = s\left(n \frac{T_s}{N}\right) = \sqrt{N} \text{IDFT}\{\mathbf{S}\}[n]$$
$$0 \leq n < N$$

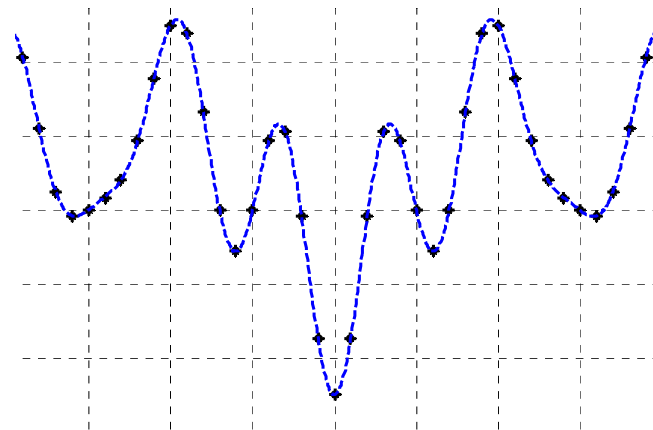
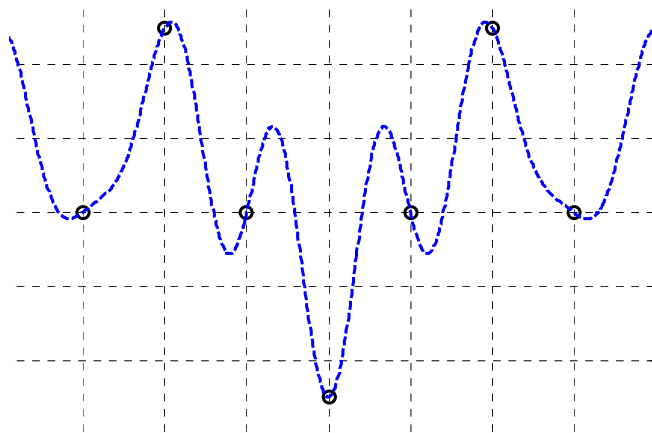


$LN$  points

$$s^{(L)}[n] = s\left(n \frac{T_s}{LN}\right) = L\sqrt{N} \text{IDFT}\{\tilde{\mathbf{S}}\}[n]$$
$$0 \leq n < LN$$

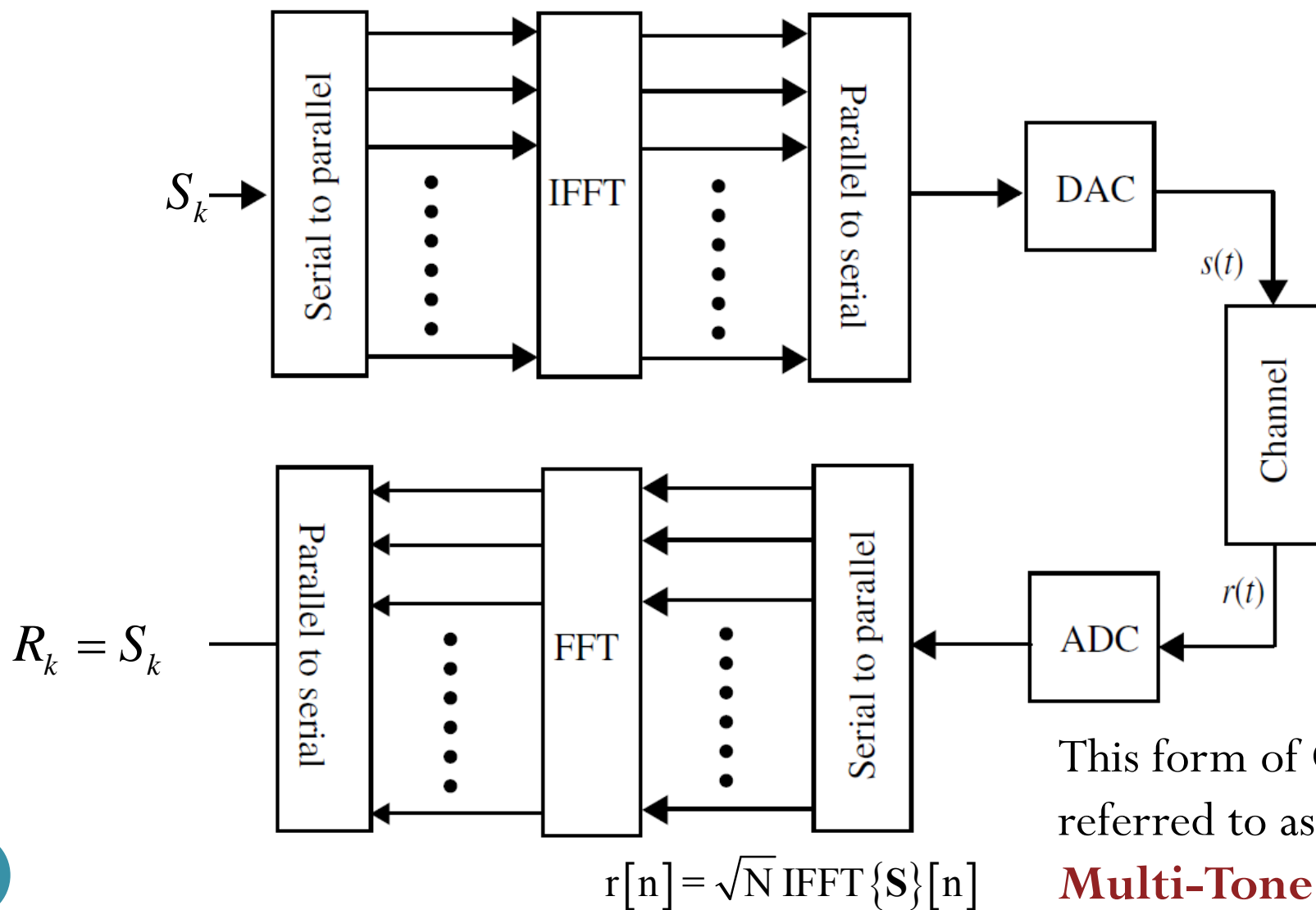
Zero padding:

$$\tilde{S}_k = \begin{cases} S_k, & 0 \leq k < N \\ 0, & N \leq k < LN \end{cases}$$



# OFDM implementation by IFFT/FFT

$$s^{(L)}[n] = s\left(n \frac{T_s}{LN}\right) = L\sqrt{N} \text{IFFT}^{(L)}\{\tilde{\mathbf{S}}\}[n]$$



This form of OFDM is often referred to as **Discrete Multi-Tone (DMT)**.

# OFDM with Memoryless Channel

$$h(t) = \beta\delta(t)$$

[should be  $h(t) = \beta\delta(t - \tau)$ ]

$$r(t) = h(t) * s(t) + w(t) = \beta s(t) + w(t)$$

Additive white Gaussian noise

Sample every  $T_s/N$

$$r[n] = \beta s[n] + w[n]$$

$$s[n] = \sqrt{N} \text{IFFT}\{S\}[n]$$

FFT

$$R_k = \frac{1}{\sqrt{N}} \text{FFT}\{\mathbf{r}\}[n] = \beta S_k + \frac{1}{\sqrt{N}} W_k$$

Sub-channel are independent.

(No ICI)

# Channel with Finite Memory

Discrete time baseband model:

$$y[n] = \{h\}^* s[n] + w[n] = \sum_{m=0}^{\nu} h[m] s[n-m] + w[n]$$

[Tse Viswanath, 2005, Sec. 2.2.3]

where  $h[n] = 0$  for  $n < 0$  and  $n > \nu$

$$w[n] \stackrel{i.i.d.}{\sim} \mathcal{CN}(0, N_0)$$

We will assume that  $\nu \ll N$

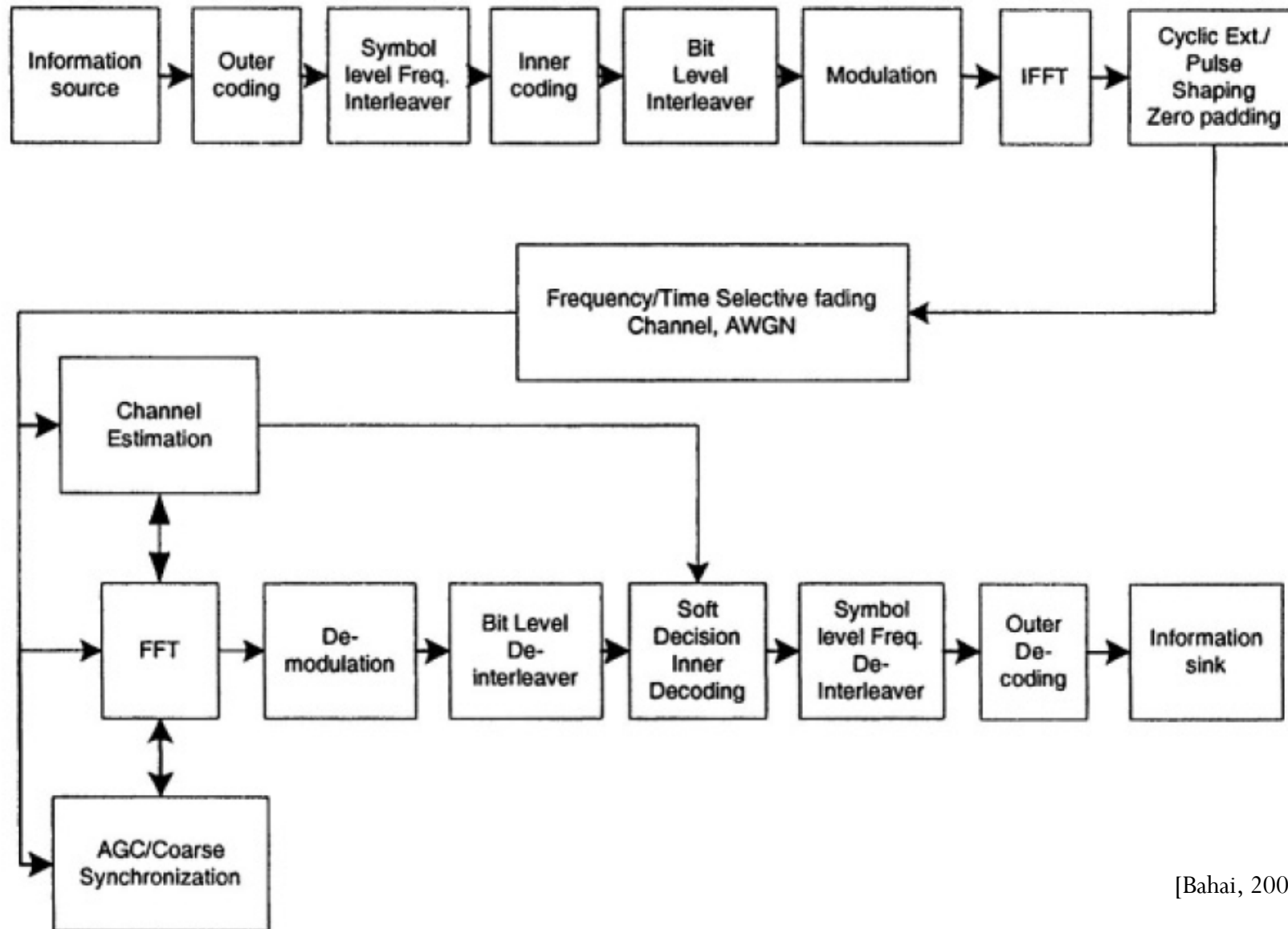
Remarks:

$Z = X + jY$  is a **complex Gaussian** if  $X$  and  $Y$  are jointly Gaussian.

If  $X, Y$  is i.i.d.  $\mathcal{N}(0, \sigma^2)$ , then  $Z = X + iY \sim \mathcal{CN}(0, \sigma_Z^2)$  where  $\sigma_Z^2 = 2\sigma^2$  with

$$f_Z(z) = f_{X,Y}(\operatorname{Re}\{z\}, \operatorname{Im}\{z\}) = \frac{1}{\pi\sigma_Z^2} e^{-\frac{|z|^2}{\sigma_Z^2}}.$$

# OFDM Architecture



[Bahai, 2002, Fig 1.11]